9 Interoperability Aspects of Fuzzy Sets
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Fuzzy sets and a family of $\alpha$ -cuts

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The representation theorem offers an important insight into links between a given fuzzy set and its $\alpha$-cuts.

In essence, any fuzzy set can be represented as an infinite family of $\alpha$-cuts:

$$A = \bigcup_{0<\alpha} \alpha A_{\alpha}$$
From fuzzy set to a family of sets

$$x \in x_{\alpha}$$

reconstruction
From fuzzy set to a family of sets

reconstruction

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From fuzzy set to a family of sets
An optimization

Is there an optimal level $a$ that optimizes a single $\alpha$-cut of $A$ so that $A_\alpha$ approximates $A$ to the highest extent?

Performance index

$$Q = \int_{x \notin A_\alpha} A(x)dx + \int_{x \in A_\alpha} (1 - A(x))dx$$

$$\min_{\alpha} Q = Q(\alpha_{\text{opt}}) \text{ viz. } \alpha_{\text{opt}} = \arg \min_{\alpha} Q(\alpha)$$

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Triangular fuzzy sets - optimization

A(x) = \text{max} \ (1-x/b, \ 0) \text{ defined for positive values of } x

Performance index

\[ Q = \int_{b(1-\alpha)}^{b} \left(1 - \frac{x}{b}\right)dx + \int_{0}^{b(1-\alpha)} \left(1 - 1 + \frac{x}{b}\right)dx \]

After integration

\[ Q = b - b(1-\alpha) + b(1-\alpha)^2 - b/2 \]

The minimum of Q, \( \frac{dQ}{d\alpha} = 0 \) is attained for \( \alpha = \frac{1}{2} \).
Set-based approximation of fuzzy sets

By approximating fuzzy sets by a finite family of sets we can directly exploit well-developed techniques of interval analysis and combine the partial results into a single fuzzy set (result).
Fuzzy sets and interfaces

Fuzzy sets do not exist in real-world (sets do not as well).

To interact with the world one has to construct interfaces (encoders and decoders)
Fuzzy sets and interfaces

The need for building interfaces exists in case of sets (interval analysis); here we encounter well-known constructs of analog-to-digital (AD) and digital-to-analog (DA) converters.
Fuzzy sets and interfaces

There are two functional modules:

**Encoders** The objective is to translate input data into some internal format acceptable for processing at level of fuzzy sets

**Decoders** The objective is to convert the results of processing of fuzzy sets into some format acceptable by the external world (typically in the form of some numeric quantities)

For encoding and decoding we engage a collection of fuzzy sets – information granules
Given is a collection of fuzzy sets $A_1, A_2, \ldots, A_c$; express some numeric input $x$ in $\mathbb{R}$ in terms of these fuzzy sets

$$x \rightarrow [A_1(x) \ A_2(x) \ldots \ A_c(x)]$$

*Nonlinear* mapping from $\mathbb{R}$ to $c$-dimensional unit hypercube
Decoding process

(a) decoding completed on a basis of a single fuzzy set

(b) Decoding realized on a basis of a certain finite family of fuzzy sets and levels of their activation.
Decoding process: a single fuzzy set

Single fuzzy set $B \rightarrow$ develop a single numeric representative
Single fuzzy set decoding: main strategies

**Mean of maxima.** Determine the arguments of $X$ for which this membership function achieves its maximal values. Denote them by $\bar{x}_1, \bar{x}_2, \ldots, \bar{x}_p$. The decoding is taken as the average of these values,

$$\hat{x} = \frac{\bar{x}_1 + \bar{x}_2 + \ldots + \bar{x}_p}{p}$$

**Centre of area.** Determine $\hat{x}$ such that it results in the equal areas below the membership function positioned on the left and on the right from this representative

$$\int_{-\infty}^{\hat{x}} B(x)dx = \int_{\hat{x}}^{\infty} B(x)dx$$

**Centre of gravity** Here the result of decoding is computed as follows

$$\hat{x} = \frac{\int B(x)dx}{\int_{\hat{x}}^{\infty} B(x)dx}$$
Single fuzzy set decoding: centre of gravity

Solution to the following optimization problem

\[ V = \int_{\hat{x}} B(x)[x - \hat{x}]^2 \, dx \]

\[ \frac{dV}{d\hat{x}} = 0 \]

\[ 2 \int_{\hat{x}} B(x)[x - \hat{x}] \, dx = 0 \]
Single fuzzy set decoding: augmented strategies

Augmented centre of gravity

\[ \hat{x} = \frac{\int_{x \in X: B(x) \geq \beta} B(x) dx}{\int_{x \in X: B(x) \geq \beta} B(x) dx} \]

\[ \hat{x} = \frac{\int_{x \in X: B(x) \geq \beta} B^\gamma(x) dx}{\int_{x \in X: B(x) \geq \beta} B^\gamma(x) dx} \]

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Single fuzzy set decoding: some general requirements

Requirements implied by:

• Monotonicity with respect to changeable membership functions

• Graphically motivated requirements (symmetry, translation, scaling…)

• The use of logic operations and logic modifiers
Decoding: a collection of fuzzy sets

- One-dimensional case
- Multivariable case

ENCODER

DECODER

Numeric Input (multidimensional)

Granular representation

Numeric Output (multidimensional)
Decoding: one-dimensional (scalar) case

Codebook – a finite family of fuzzy sets \{A_1, A_2, \ldots, A_c\}
One-dimensional decoding

Assume

(a) the fuzzy sets of the codebook \{A_i\}, i=1, 2,...,c form a fuzzy partition, \(\sum_{i=1}^{c} A_i(x) = 1\), and for each \(x\) in \(X\) at least one element of the codebook is “activated”, that is \(A_i(x) > 0\)

(b) for each \(x\) only two neighboring elements of the codebook are “activated” that is \(A_1(x) = 0, ..., A_{i-1}(x) = 0, A_i(x) > 0, A_{i+1}(x) > 0, A_{i+2}(x) = ... = A_c(x) = 0\)

(c) the decoding is realized as a weighted sum of the activation levels and the prototypes of the fuzzy sets \(v_i\), namely \(\hat{x} = \sum_{i=1}^{c} A_i(x)v_i\)

Then the elements of the codebook described by piecewise linear membership functions

\[
A_i(x) = \begin{cases} 
\frac{x - v_{i-1}}{v_i - v_{i-1}} & \text{if } x \in [v_{i-1}, v_i] \\
\frac{x - v_{i+1}}{v_i - v_{i+1}} & \text{if } x \in [v_i, v_{i+1}] \\
v_i - v_{i+1} & \text{if } x \in [v_{i+1}, v_{i}] 
\end{cases}
\]

lead to the zero decoding error (lossless compression) meaning that \(\hat{x} = x\).
Multivariable encoding and decoding: a global picture

- **Encoding**
  - Prototypes $v_1, v_2, \ldots, v_c$

- **Decoding**
  - Prototypes $v_1, v_2, \ldots, v_c$

**VQ**

- Use of sets – **Vector Quantization (VQ)**

**FVQ**

- Use of fuzzy sets – **Fuzzy Vector Quantization (FVQ)**
The codebook formed through fuzzy clustering (FCM) producing a finite collection of prototypes $v_1, v_2, \ldots, v_c$.

Given any new input $x$ we realize its encoding and decoding.

Let us recall

**Encoding** – representation of $x$ in terms of the prototypes

**Decoding** – development of external representation of the result of processing realized at the level of information granules
The optimization problem

\[
\sum_{i=1}^{c} u_i^m \| x - v_i \|^2
\]

Minimize w.r.t. \( u_i \) subject to

\[
u_i(x) \in [0,1], \quad \sum_{i=1}^{c} u_i(x) = 1
\]

\[
u_i(x) = \frac{1}{\sum_{j=1}^{c} \left( \frac{\| x - v_i \|}{\| x - v_j \|} \right)^{2/(m-1)}}
\]
Fuzzy Vector Quantization: Decoding

Reconstruct original multidimensional input $x$

$$u_i(x) \in [0,1], \quad \sum_{i=1}^{c} u_i(x) = 1$$

$$Q_2(\hat{x}) = \sum_{i=1}^{c} (\hat{x} - v_i)^2$$

minimize

$$\hat{x} = \frac{\sum_{i=1}^{c} u_i^m v_i}{\sum_{i=1}^{c} u_i^m}$$

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Fuzzy Vector Quantization: Decoding error

\[ u_i(x) \in [0,1], \quad \sum_{i=1}^{c} u_i(x) = 1 \]
Fuzzy encoding and decoding with possibility and necessity measures

Consider a family of fuzzy sets $A_1, A_2, \ldots, A_c$

Input datum $X$ could be either a fuzzy set or a numeric quantity

\[
Poss(A_i, X) = \sup_{x \in X} [X(x) \cdot t_{A_i}(x)]
\]

\[
Nec(A_i, X) = \inf_{x \in X} [(1 - A_i(x)) \cdot s_X(x)]
\]
Possibility and necessity encoding

\[ u_i(x) \in [0,1], \quad \sum_{i=1}^{\mathcal{X}} u_i(x) = 1 \]

\[ \sum_{i=1}^{\mathcal{X}} \in \]

\[ = \]

\[ c_i \]

\[ \sum_{i=1}^{\mathcal{X}} u_i(x) = 1 \]

\[ = \]

\[ c_i \]

\[ \sum_{i=1}^{\mathcal{X}} \in \]

\[ = \]

\[ c_i \]

\[ \sum_{i=1}^{\mathcal{X}} u_i(x) = 1 \]

\[ = \]

\[ c_i \]

\[ \sum_{i=1}^{\mathcal{X}} u_i(x) = 1 \]

\[ = \]

\[ c_i \]

\[ \sum_{i=1}^{\mathcal{X}} u_i(x) = 1 \]

\[ = \]

\[ c_i \]

\[ \sum_{i=1}^{\mathcal{X}} ]
Possibility and necessity encoding: example

\[ \sum_{i=1}^{c} u_i(x) = 1 \]

\[ u_i(x) \in [0,1], \quad \sum_{i=1}^{c} u_i(x) = 1 \]

\[ X = [0.0 \ 0.2 \ 0.8 \ 1.0 \ 0.9 \ 0.5 \ 0.1 \ 0.0] \]

\[ A_i = [0.6 \ 0.5 \ 0.4 \ 0.5 \ 0.6 \ 0.9 \ 1.0 \ 1.0] \]

\[ \text{Poss}(A_i, X) = \max (0.0, 0.5, 0.4, 0.5, 0.6, 0.5, 0.1, 0.0) = 0.6 \]

\[ \text{Nec}(A_i, X) = \min(0.4, 0.5, 0.8, 1.0, 0.9, 0.5, 0.1, 0.0) = 0.0 \]
Encoding and decoding: an overview

$u_i(x) \in [0,1], \quad \sum_{i} u_i(x) = 1$
Decoding

Given the nature of encoding (possibility and necessity measures), the decoding is regarded as a certain “inverse” problem in terms of fuzzy relational equations:

Possibility measure – sup-t composition

Necessity measure – inf-s composition
Decoding – possibility measure

Possibility measure – sup-t composition

\[ \hat{X}(x) = A(x) \rightarrow \lambda = \begin{cases} 1 & \text{if } A(x) \leq \lambda \\ \lambda & \text{otherwise} \end{cases} \]

\[ \hat{X}(x) = A(x) \rightarrow \lambda = \sup[a \in [0,1] | a A(x) \leq \lambda]\]

\[ \hat{X} = \bigcap_{i=1}^{c} \hat{X}_i \]
Decoding – necessity measure

Necessity measure – inf-s composition

\[ \tilde{X}(x) = (1 - A(x)) \epsilon \mu = \begin{cases} \mu, & \text{if } 1 - A(x) < \mu \\ 0, & \text{otherwise} \end{cases} \]

\[ \tilde{X}(x) = (1 - A_i(x)) \epsilon \mu = \inf \{ a \in [0,1] | as(1 - A(x)) \geq \mu \} \]

\[ \tilde{X} = \bigcup_{i=1}^{c} \tilde{X}_i \]

\[ \tilde{X} \subseteq X \subseteq \hat{X} \]
Decoding: example

possibility measure
Decoding: example

necessity measure

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Decoding: example

bounds of possibility and necessity measure
Taxonomy of data structure with the use of shadowed sets

Core structure

Shadowed data structure

Uncertain data structure
Core data structure patterns that belong to a core of at least one or more shadowed sets

Core data structure = \{ x | \exists_i x \in \text{core}(A_i) \}
Shadowed data structure This structure is formed by patterns that do not belong to core of any of the shadowed sets but fall within the shadow of one or more shadowed sets.

\[
\text{Shadowed data structure} = \{ \mathbf{x} \mid \exists \mathbf{x} \in \text{shadow}(A_i) \text{ and } \forall_i \mathbf{x} \notin \text{core}(A_i) \}\]
Uncertain data structure The patterns belonging to this structure are those that are left out from all shadows

Uncertain data structure = \{ x \mid \forall_i x \not\in \text{shadow}(A_i) \text{ and } \forall_i x \not\in \text{core}(A_i) \}
Three-valued characterization of data structure with shadowed sets
Three-valued characterization of data structure: example
Three-valued characterization of data structure: example

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